# Reciprocity laws and $\Delta_{0}$-definability 

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Plan
-What is $\Delta_{0}$-definability?

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- Reciprocity schema


## Plan

-What is $\Delta_{0}$-definability?

- Dedekind sums are $\Delta_{0}$-definable
- Rademacher-Dieter-Knuth-Dedekind sums are $\Delta_{0}$-definable
- Reciprocity schema
- Conclusion


## What is $\Delta_{0}$-definability?

$x$ is not prime nor 0 nor 1
$(\exists u) \quad(\exists v) \quad(x=u v) \wedge(u \neq x) \wedge(u \neq x)$

## What is $\Delta_{0}$-definability?

$x$ is not prime nor 0 nor 1

$$
(\exists u)_{u<x}(\exists v)_{v<x}(x=u v) \wedge(u \neq x) \wedge(v \neq x)
$$

## What is $\Delta_{0}$-definability?

Major open problem

Find a "simple" artithmetical relation
NOT $\Delta_{0}$-definable

## What is $\Delta_{0}$-definability?

## Exemple

$$
z=x^{y}
$$

## What is $\Delta_{0}$-definability?

## Exemple

$$
z=x^{y}
$$

$$
y=2 \wedge z=x \cdot x \text { and } y=3 \wedge z=x \cdot x \cdot x \text { and } \ldots
$$

## What is $\Delta_{0}$-definability?

## Exemple

$$
z=x^{x}
$$

IS $\Delta_{0}$-definable

## What is $\Delta_{0}$-definability?

Open exemple
$z$ is the $n$-th prime number
IS NOT KNOWN TO BE $\Delta_{0}$-definable

## A new $\Delta_{0}$-definable relation and a standard method

Let $d$ and $c$ be integers, $c \neq 0$. The Dedekind sum is defined by

$$
s(d, c)=\sum_{k=1}^{k=|c|}\left(\left(\frac{k}{c}\right)\right)\left(\left(\frac{k d}{c}\right)\right)
$$

where $((x))=0$ if $x \in Z$, else $x-[x]-\frac{1}{2}$.

# A new $\Delta_{0}$-definable relation and a standard method 

$s(d, c)$ is a rational number, but
$12 c \times s(d, c)$ is a rational integer.

## A new $\Delta_{0}$-definable relation and a standard method

$s(d, c)$ is a rational number, but
$12 c \times s(d, c)$ is a rational integer.

$$
s(d, c)=\frac{c-1}{12 c}(4 c d-2 d-3 c)-\frac{1}{c} \sum_{n=1}^{c-1} n\left\lfloor\frac{d n}{c}\right\rfloor
$$

A new $\Delta_{0}$-definable relation and a standard method

Theorem : $12 c \times s(d, c)$ is $\Delta_{0}$-definable.

## A new $\Delta_{0}$-definable relation and a standard method

Lemma (Ph. Barkan, 1977)

$$
s(d, c)=\frac{1}{12}\left(-3+\frac{d+d^{-1} \bmod c}{c}-\sum_{i=1}^{i=r}(-1)^{i} a_{i}\right)
$$

where $\left(a_{i}\right)_{0 \leq i \leq r}$ is the sequence of the continued fraction development

$$
\frac{d}{c}=a_{0}+\frac{1}{a_{1}+\frac{1}{a_{2}+\ldots+\frac{1}{a_{r}}}}
$$

## A new $\Delta_{0}$-definable relation and a standard method

Lemma (Woods, 1981)
$z=\sum_{i=0}^{i=g(\mathbf{x})} f(i, \mathbf{x})$ is $\Delta_{0}$-definable provided

- the graphs of $f(i, \mathbf{x})$ and $g(\mathbf{x})$ are $\Delta_{0}$ - definable
- $f$ is polynomially bounded
- $g$ is polylogarithmicaly bounded (i.e. there exists a polynomial function $\psi$ with positive integer coefficients, such that $\left.g(\mathbf{x}) \leq \psi\left(\left\lfloor\log _{2}\left(x_{1}\right)\right\rfloor, \ldots,\left\lfloor\log _{2}\left(x_{k}\right)\right\rfloor\right)\right)$


## A new $\Delta_{0}$-definable relation and a standard method

Lemma (H.-A. E. 2010)
Let $c$ and $d$ be two positive integers. The sequence of the coefficients $\left(a_{i}(d, c)\right)_{0 \leq i \leq r(d, c)}$ of the continued fraction development of $\frac{d}{c}$ is $\Delta_{0}$-definable.

Generalization of Dedekind sum

Main result

## Generalization of Dedekind sum

Let us consider now the following generalization known (?) as the Rademacher-Dieter-Knuth-Dedekind sum :

$$
r_{u}(d, c)=\sum_{k=1}^{k=|c|}\left(\left(\frac{k}{c}\right)\right)\left(\left(\frac{k d+u}{c}\right)\right)
$$

where $u \in Z$.

## Generalization of Dedekind sum

$r_{u}(d, c)$ is a rational number, but

$$
12 c \times r_{u}(d, c) \text { is a rational integer. }
$$

Theorem : $12 c \times r_{u}(d, c)$ is $\Delta_{0}$-definable.

## Generalization of Dedekind sum

Reciprocity law (U. Dieter, 1959)
$r_{u}(d, c)+r_{u}(c, d)=\frac{1}{12}\left(\frac{d}{c}+\frac{c}{d}+\frac{1+6\lfloor u\rfloor\lceil u\rceil}{c d}-6\left\lfloor\frac{u}{c}\right\rfloor-3 e(d, u)\right)$
for $d>u>0$ and $d \geq c>0$ and $e(d, u)=0$ if $u>0$ and $u \equiv 0 \bmod c$, else $e(c, u)=1$.

## Generalization of Dedekind sum

$$
r_{0}=c \quad r_{1}=d \quad u_{0}=u
$$

Euclidean algorithm
$r_{i+2}=r_{i} \bmod r_{i+1}$
$u_{i+1}=u_{i} \bmod r_{i+1}$
$\alpha_{i}=\left\lfloor\frac{r_{i}}{r_{i+1}}\right\rfloor$
$\beta_{i}=\left\lfloor\frac{u_{i}}{r_{i+1}}\right\rfloor$
$r_{l-1}=\alpha_{l-1} r_{l}$
$u_{I-1}=\beta_{I-1} r_{I}+u_{I}$

## Generalization of Dedekind sum

$$
\begin{array}{llll}
p_{0}=a_{0} & p_{1}=a_{0} a_{1}+1 & q_{0}=1 & q_{1}=a_{1} \\
\alpha_{i} p_{i-1}+p_{i-2}=p_{i} & \alpha_{i} q_{i-1}+q_{i-2}=q_{i} &
\end{array}
$$

## Generalization of Dedekind sum

Algorithm (D. E. Knuth, 1977)

$$
\begin{gathered}
12 c \times r_{u}(d, c)=c\left(\sum_{j=0}^{j=l-1}(-1)^{j}\left(\alpha_{j}-6 \beta_{j}-3 e\left(r_{j+1}, u_{j}\right)\right)\right) \\
+\left(d+(-1)^{l-1} p_{l}+6 \sum_{j=0}^{j=l-1}(-1)^{j} \beta_{j}\left(u_{j}+u_{j+1}\right) p_{j+1}\right)
\end{gathered}
$$

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\end{gathered}
$$

## Generalization of Dedekind sum

The idea is to code $\left(\beta_{j}\right)_{0 \leq j \leq 1}$ with $\left(\Gamma_{1}, \Gamma_{2}\right)$ such that

$$
\frac{\Gamma_{1}}{\Gamma_{2}}=\beta_{0}+\frac{1}{\beta_{1}+\frac{\ldots}{\cdots+\frac{1}{\beta_{l-1}}}}
$$

## Generalization of Dedekind sum

The idea is to code $\left(\beta_{j}\right)_{0 \leq j \leq I}$ with $\left(\Gamma_{1}, \Gamma_{2}\right)$ such that

$$
\frac{\Gamma_{1}}{\Gamma_{2}}=\beta_{0}+\frac{1}{\beta_{1}+\frac{\ldots}{\ldots+\frac{1}{\beta_{l-1}}}}
$$

The code is polynomially bounded :
$\Gamma_{1}<u d$ and $\Gamma_{2}<d$.

## Generalization of Dedekind sum

The idea is to code $\left(\beta_{j}\right)_{0 \leq j \leq I}$ with $\left(\Gamma_{1}, \Gamma_{2}\right)$ such that

$$
\frac{\Gamma_{1}}{\Gamma_{2}}=\beta_{0}+\frac{1}{\beta_{1}+\frac{\ldots}{\ldots+\frac{1}{\beta_{l-1}}}}
$$

The decoding function $\alpha_{j}\left(\Gamma_{1}, \Gamma_{2}\right)$ is rudimentary.

## Generalization of Dedekind sum

$z=\beta_{j}$ is definned by

$$
\begin{gathered}
\exists\left(\Gamma_{1}\right)_{<u d} \exists\left(\Gamma_{2}\right)_{<d} \exists\left(\Gamma_{3}\right)_{<u}\left(z=\alpha_{j}\left(\Gamma_{1}, \Gamma_{2}\right)\right) \wedge \\
\left(\forall i\left(\Gamma_{3}+\sum_{k=l-1}^{k=i} \alpha_{j}(u, v) r_{i+1}<r_{i}\right)\right)
\end{gathered}
$$

and simillary for $z=u_{j}$.

## Reciprocity laws

RL schema :

$$
\left\{\begin{array}{l}
f(a, b, \vec{x})=f(a-b, b, \vec{x}) \text { if } \mathrm{a}>\mathrm{b} \\
f(a, 0, \vec{x})=g(a, \vec{x}) \\
f(a, b, \vec{x})=h(f(b, a, \vec{x}), a, b, \vec{x}) \\
f(a, b, \vec{x}) \leq p o l y(a, b, \vec{x}) ? ?
\end{array}\right.
$$

defines $f$ from $g$ and $h$ with N as domain and codomain.

## Reciprocity laws

summing RL schema :

$$
\left\{\begin{array}{l}
f(a, b, \vec{x})=f(a-b, b, \vec{x}) \text { if } \mathrm{a}>\mathrm{b} \\
f(a, 0, \vec{x})=g(a, \vec{x}) \\
f(a, b, \vec{x})+f(b, a, \vec{x})=h(a, b, \vec{x}) \\
\mathrm{h}(a, b, \vec{x}) \leq \operatorname{poly}(a, b, \vec{x})
\end{array}\right.
$$

defines $f$ from $g$ and $h$ with N as domain and codomain.

## Reciprocity laws

Proposition : The set of rudimentary functions is closed under the RL summing schema.

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Proof:
$f(a, b, \vec{x})=\sum_{j=1}^{j=1}(-1)^{j+1} h\left(r_{j+1}, r_{j}, \vec{x}\right)+(-1)^{\prime} g\left(r_{l}, \vec{x}\right)$
$r_{0}=a \quad r_{1}=b \quad$ Euclidean algorithm
$r_{i+2}=r_{j} \bmod r_{i+1} \quad r_{l-1}=\alpha_{l-1} r_{l}$
$r_{l}=\operatorname{gcd}(a, b)$

## Reciprocity laws

|  | Recip. Law | log. long <br> pol. bounded <br> rec. schema | classical <br> pol bounded <br> rec. schema |
| :---: | :---: | :---: | :---: |
| complete | $?$ | $?$ | $?$ |
| sum weakened | closed | closed | equiv. counting |

## Reciprocity laws

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| :---: | :---: | :---: | :---: |
| complete | $\leftarrow$ | $\leftarrow$ | $\leftarrow$ |
| sum weakened | closed | closed | equiv. counting |

Conclusion and further work

More generalization of Dedekind sum

$$
s(a, b, c)=\sum_{k=1}^{k=c}\left(\left(\frac{a k}{c}\right)\right)\left(\left(\frac{b k}{c}\right)\right)
$$

Conclusion and further work

Reciprocity (Rademacher, 1954)

$$
s(a, b, c)+s(b, c, a)+s(c, a, b)=\frac{1}{4}-\frac{1}{12}\left(\frac{c}{a b}+\frac{a}{b c}+\frac{b}{a c}\right)
$$

Conclusion and further work

Conjecture : it is $\Delta_{0}^{\sharp}$ definable

