Reciprocity laws and Δ_0 -definability

Henri-Alex Esbelin

Clermont-Ferrand Universities

JAF 32, Athens

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

• What is Δ_0 -definability?

イロト イロト イモト イモト 一日

- What is Δ_0 -definability?
- Dedekind sums are Δ_0 -definable

ヘロト ヘロト ヘビト ヘビト 一日

- What is Δ_0 -definability?
- Dedekind sums are Δ_0 -definable
- Rademacher-Dieter-Knuth-Dedekind sums are Δ_0 -definable

- What is Δ_0 -definability?
- Dedekind sums are Δ_0 -definable
- Rademacher-Dieter-Knuth-Dedekind sums are Δ_0 -definable

• Reciprocity schema

- What is Δ_0 -definability?
- Dedekind sums are Δ_0 -definable
- Rademacher-Dieter-Knuth-Dedekind sums are Δ_0 -definable

- Reciprocity schema
- Conclusion

x is not prime nor 0 nor 1

$$(\exists u)$$
 $(\exists v)$ $(x = uv) \land (u \neq x) \land (u \neq x)$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへで

x is not prime nor 0 nor 1

$$(\exists u)_{u < x} (\exists v)_{v < x} (x = uv) \land (u \neq x) \land (v \neq x)$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへで

Major open problem

Find a "simple" artithmetical relation

NOT Δ_0 -definable

< □ > < □ > < 三 > < 三 > < 三 > < 三 > < ○ < ○

Exemple

$$z = x^y$$

・ロト・西ト・ヨト・ヨト・日下 ひゃつ

Exemple

$$z = x^y$$

$$y = 2 \land z = x.x$$
 and $y = 3 \land z = x.x.x$ and ...

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

Exemple

$$z = x^{x}$$

IS Δ_0 -definable

うせん 山田 (山田) (山) (山)

Open exemple

z is the n-th prime number

IS NOT KNOWN TO BE Δ_0 -definable

Let *d* and *c* be integers, $c \neq 0$. The Dedekind sum is defined by

$$s(d,c) = \sum_{k=1}^{k=|c|} \left(\left(\frac{k}{c} \right) \right) \left(\left(\frac{kd}{c} \right) \right)$$

where ((x)) = 0 if $x \in Z$, else $x - [x] - \frac{1}{2}$.

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへで

s(d, c) is a rational number, but

 $12c \times s(d, c)$ is a rational integer.



s(d, c) is a rational number, but

 $12c \times s(d, c)$ is a rational integer.

$$s(d,c) = \frac{c-1}{12c}(4cd - 2d - 3c) - \frac{1}{c}\sum_{n=1}^{c-1}n\left\lfloor\frac{dn}{c}\right\rfloor$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへで

Theorem : $12c \times s(d, c)$ is Δ_0 -definable.



Lemma (Ph. Barkan, 1977)

$$s(d, c) = \frac{1}{12} \left(-3 + \frac{d + d^{-1} \mod c}{c} - \sum_{i=1}^{i=r} (-1)^i a_i \right)$$

where $(a_i)_{0 \le i \le r}$ is the sequence of the continued fraction development

$$\frac{d}{c} = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{\dots}{\dots + \frac{1}{a_r}}}}$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへで

Lemma (Woods, 1981)

$$z = \sum_{i=0}^{i=g(\mathbf{x})} f(i, \mathbf{x})$$
 is Δ_0 -definable provided

- the graphs of $f(i, \mathbf{x})$ and $g(\mathbf{x})$ are Δ_0 definable
- f is polynomially bounded

• *g* is polylogarithmically bounded (i.e. there exists a polynomial function ψ with positive integer coefficients, such that $g(\mathbf{x}) \leq \psi(\lfloor \log_2(x_1) \rfloor, ..., \lfloor \log_2(x_k) \rfloor))$

Lemma (H.-A. E. 2010)

Let *c* and *d* be two positive integers. The sequence of the coefficients $(a_i(d, c))_{0 \le i \le r(d,c)}$ of the continued fraction development of $\frac{d}{c}$ is Δ_0 -definable.

Main result

▲□▶ ▲□▶ ▲目▶ ▲目▶ ▲□ ● ● ●

Let us consider now the following generalization known (?) as the Rademacher-Dieter-Knuth-Dedekind sum :

$$r_u(d,c) = \sum_{k=1}^{k=|c|} \left(\left(\frac{k}{c} \right) \right) \left(\left(\frac{kd+u}{c} \right) \right)$$

where $u \in Z$.

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

 $r_u(d, c)$ is a rational number, but

 $12c \times r_u(d, c)$ is a rational integer.

Theorem : $12c \times r_u(d, c)$ is Δ_0 -definable.

Reciprocity law (U. Dieter, 1959)

$$r_u(d,c)+r_u(c,d)=\frac{1}{12}\left(\frac{d}{c}+\frac{c}{d}+\frac{1+6\lfloor u\rfloor\lceil u\rceil}{cd}-6\lfloor\frac{u}{c}\rfloor-3e(d,u)\right)$$

for d > u > 0 and $d \ge c > 0$ and e(d, u) = 0 if u > 0 and $u \equiv 0 \mod c$, else e(c, u) = 1.

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへ⊙

$$r_0 = c$$
 $r_1 = d$ $u_0 = u$

Euclidean algorithm

 $r_{i+2} = r_i \mod r_{i+1} \qquad \qquad u_{i+1} = u_i \mod r_{i+1}$ $\alpha_i = \lfloor \frac{r_i}{r_{i+1}} \rfloor \qquad \qquad \beta_i = \lfloor \frac{u_i}{r_{i+1}} \rfloor$ $r_{l-1} = \alpha_{l-1}r_l \qquad \qquad u_{l-1} = \beta_{l-1}r_l + u_l$

 $p_0 = a_0$ $p_1 = a_0 a_1 + 1$ $q_0 = 1$ $q_1 = a_1$ $\alpha_i p_{i-1} + p_{i-2} = p_i$ $\alpha_i q_{i-1} + q_{i-2} = q_i$

(日)

Algorithm (D. E. Knuth, 1977)

$$12c \times r_u(d,c) = c\left(\sum_{j=0}^{j=l-1} (-1)^j \left(\alpha_j - 6\beta_j - 3e(r_{j+1},u_j)\right)\right)$$

+
$$\left(d + (-1)^{l-1}p_l + 6\sum_{j=0}^{j=l-1}(-1)^j\beta_j(u_j + u_{j+1})p_{j+1}\right)$$

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

$$12c \times r_u(d,c) = c\left(\sum_{j=0}^{j=l-1} (-1)^j \left(\alpha_j - 6\beta_j - 3e(r_{j+1},u_j)\right)\right)$$

+
$$\left(d + (-1)^{l-1}p_l + 6\sum_{j=0}^{j=l-1} (-1)^j \beta_j (u_j + u_{j+1})p_{j+1}\right)$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 − のへで

$$12c \times r_u(d,c) = c\left(\sum_{j=0}^{j=l-1} (-1)^j \left(\alpha_j - 6\beta_j - 3e(r_{j+1}, u_j)\right)\right)$$

+
$$\left(d + (-1)^{l-1}p_l + 6\sum_{j=0}^{j=l-1} (-1)^j \beta_j (u_j + u_{j+1})p_{j+1}\right)$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 − のへで

The idea is to code $(\beta_j)_{0 \le j \le l}$ with (Γ_1, Γ_2) such that

$$\frac{\Gamma_1}{\Gamma_2} = \beta_0 + \frac{1}{\beta_1 + \frac{\cdots}{\cdots + \frac{1}{\beta_{l-1}}}}$$

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

The idea is to code $(\beta_j)_{0 \le j \le l}$ with (Γ_1, Γ_2) such that

$$\frac{\Gamma_1}{\Gamma_2} = \beta_0 + \frac{1}{\beta_1 + \frac{\cdots}{\cdots + \frac{1}{\beta_{l-1}}}}$$

The code is polynomially bounded : $\Gamma_1 < ud$ and $\Gamma_2 < d$.

The idea is to code $(\beta_j)_{0 \le j \le l}$ with (Γ_1, Γ_2) such that

$$\frac{\Gamma_1}{\Gamma_2} = \beta_0 + \frac{1}{\beta_1 + \frac{\cdots}{\cdots + \frac{1}{\beta_{l-1}}}}$$

The decoding function $\alpha_j(\Gamma_1, \Gamma_2)$ is rudimentary.

 $z = \beta_j$ is definned by

$$\exists (\Gamma_1)_{< ud} \exists (\Gamma_2)_{< d} \exists (\Gamma_3)_{< u} (z = \alpha_j(\Gamma_1, \Gamma_2)) \land$$

$$\left(\forall i \left(\Gamma_3 + \sum_{k=l-1}^{k=i} \alpha_j(u, v) r_{l+1} < r_l \right) \right)$$

▲□▶ ▲□▶ ▲豆▶ ▲豆▶ □豆 = のへ⊙

and similary for $z = u_j$.

RL schema :

$$\begin{cases} f(a, b, \vec{x}) = f(a - b, b, \vec{x}) \text{ if } a > b \\ f(a, 0, \vec{x}) = g(a, \vec{x}) \\ f(a, b, \vec{x}) = h(f(b, a, \vec{x}), a, b, \vec{x}) \\ f(a, b, \vec{x}) \le poly(a, b, \vec{x})?? \end{cases}$$

defines *f* from *g* and *h* with N as domain and codomain.

summing RL schema :

$$\begin{cases} f(a, b, \vec{x}) = f(a - b, b, \vec{x}) \text{ if } a > b \\ f(a, 0, \vec{x}) = g(a, \vec{x}) \\ f(a, b, \vec{x}) + f(b, a, \vec{x}) = h(a, b, \vec{x}) \\ h(a, b, \vec{x}) < poly(a, b, \vec{x}) \end{cases}$$

defines *f* from *g* and *h* with N as domain and codomain.

Proposition : The set of rudimentary functions is closed under the RL summing schema.

Proposition : The set of rudimentary functions is closed under the RL summing schema.

Proof :

$$f(a, b, \vec{x}) = \sum_{j=1}^{j=l} (-1)^{j+1} h(r_{j+1}, r_j, \vec{x}) + (-1)^l g(r_l, \vec{x})$$

 $r_0 = a$ $r_1 = b$ Euclidean algorithm

 $r_{i+2} = r_j \mod r_{i+1}$ $r_{l-1} = \alpha_{l-1}r_l$

 $r_l = gcd(a, b)$

	Recip. Law schema	log. long pol. bounded rec. schema	classical pol bounded rec. schema
complete	?	?	?
sum weakened	closed	closed	equiv. counting

	Recip. Law schema	log. long pol. bounded rec. schema	classical pol bounded rec. schema
complete	<i>~</i>	← ←	<i>←</i>
sum weakened	closed	closed	equiv. counting

Conclusion and further work

More generalization of Dedekind sum

$$s(a, b, c) = \sum_{k=1}^{k=c} \left(\left(\frac{ak}{c} \right) \right) \left(\left(\frac{bk}{c} \right) \right)$$

▲□▶ ▲□▶ ▲三▶ ▲三▶ - 三 - のへで

Conclusion and further work

Reciprocity (Rademacher, 1954)

$$s(a,b,c)+s(b,c,a)+s(c,a,b)=rac{1}{4}-rac{1}{12}\left(rac{c}{ab}+rac{a}{bc}+rac{b}{ac}
ight)$$

・ロト・「「「・山下・山下・山下・山下・」

Conclusion and further work

Conjecture : it is Δ_0^{\sharp} definable

▲ロト ▲ □ ト ▲ 三 ト ▲ 三 三 つへぐ